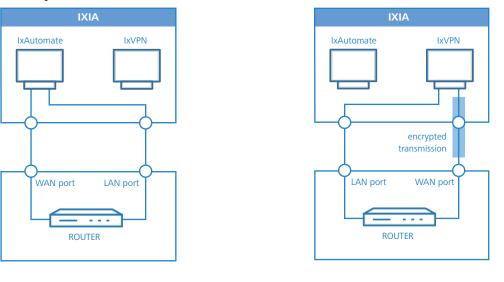
Routing performance LCOS 10.50 / 10.60

Applications for communications and entertainment are increasingly based on IP networks. In order to ensure that the necessary bandwidth performance can be provided reliably, it is important for the infrastructure's networking components to be tested thoroughly and intensively. In this techpaper, LANCOM Systems presents the methods of measuring routing and VPN performance for central site and VPN gateways as well as the respective results.

We have examined a variety of aspects for consideration when measuring the router performance. This includes transmission speeds of connections between the LAN and the Internet (WAN), and the internal data transmission in the network (LAN-LAN). Many business processes rely on secure WAN connections, which is why we have focused on determining the performance of encrypted data connections over VPN.

Test system



All of the performance values were measured in the LANCOM test laboratory. Tests were conducted with an IXIA test system. IXIA uses so-called test suites, which enable the simulation of different applications. This allows, for example, the investigation of data throughput over automatically established VPN tunnels, or the measurement of pure LAN-WAN routing performance for unidirectional and bidirectional data connections. IXIA is a leading supplier of systems which test IP-based services and infrastructures. Test systems from IXIA are employed all over the world by network component manufacturers and other organizations to help assure the functionality and reliability of complex IP networks, devices, and applications.

Figure 1: IXIA test system for routing connections and encrypted VPN connections between LAN and WAN



The measurement of data transmission itself uses either a fixed frame size or a combination of frame sizes which reflects a typical flow of data. These combinations are known as "Internet Mix", or IMIX for short. The type of IMIX which is applied significantly affects the test results because packet size has a strong influence on a connection's performance. By selecting the appropriate ports on the router being tested, it is possible to test connections between the LAN and the WAN, and also pure LAN-LAN connections.

The setup for measuring transfer rates > 1 GBit/s represents a large central-site scenario. In this scenario, several central sites can also work as a network, which is why an intermediate router with BGP ensures that the packets for each tunnel pass through the respective central site (measured values from <u>Table 11</u>).

Two scenarios are used for the vRouter measurements. One is the transmission from a PC in the LAN to a PC in another LAN or WAN by using the vRouter in between. The second scenario adds two vRouters which are connected over WAN and use a VPN tunnel for encrypted data transfer.

Firmware on the devices under test is LCOS 10.50 (unless otherwise stated).

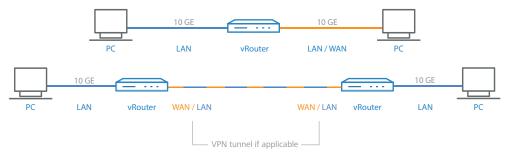


Figure 2: Schematic view of the vRouter test system scenarios



Routing performance (UDP)

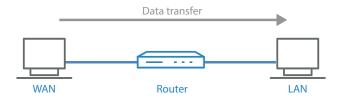


Figure 3: Schematic view of the test system

The measurement of routing performance involves the determination of the maximum data throughput which can be achieved before a router starts rejecting packets. This measurement uses UDP packets of various sizes in order to simulate the performance with different applications. Ethernet frame sizes range from 64 bytes for the smallest to 1518 bytes for the largest frames. Tests on different router models demonstrate the influence of the different hardware platforms (processor, interfaces). Measurements initially determine the frame rate, which is a good performance

indicator of the tested hardware. With normal routing, the frame rate is fairly constant even with different frame sizes. This is because only the header is inspected during routing, a process which is largely independent of the size of the frames being routed. For this reason, only the typical frame rates are given in the tables.

The throughput for a certain frame size (or even a mix of sizes, see IMIX on page 05) can be approximately calculated by multiplication with the frame rate. When the frame rate is constant, data throughput depends directly on the frame size because the larger the frames, the larger is the data volume that can be transmitted. The maximum number of frames transmitted per second is limited by the performance of the interfaces and the transmission medium.

Measurement of the routing performance relates to the size of the Ethernet frames. To compare packet sizes for particular applications, it is necessary to subtract the header. For a frame of 512 bytes, the result is a UDP datagram size of 474 bytes (512 bytes - 14 bytes Ethernet header - 4 bytes FCS trailer - 20 bytes IP header) and, after subtracting the UDP header (8 bytes), the UDP payload is 466 bytes. To investigate routing performance, in this paper two different applications are

considered.

- → For WAN-LAN routing, data received from the WAN is forwarded to a peer in the LAN.
- \rightarrow For LAN-LAN routing, data remains within the local-area network and is passed from one LAN port to another.

The measurements show that the throughput increases almost linearly with the frame size until the limit of the Gigabit interface is reached.



Routing performance (TCP)

Using UDP, the maximum performance is easily measured. However, a large portion of data transfers utilizes TCP. Therefore, it is important to take a closer look at those scenarios as well.

The TCP measurement is accomplished by using iperf, a program used to measure the data throughput between two computers. They are connected to a LANCOM router. One computer acts as a server in the WAN, sending data packets to the other, which acts as a client in the LAN, representing a typical download scenario.

The TCP measurement done with iperf did not use NAT since no direct measurement with NAT is possible using iperf. The correlation between the test results and the results seen at a WAN connection including NAT in a live environment is very high. The causes are the differences of the traffic structure of iperf compared to an HTTP download and the LCOS firewall which offers a similar performance with NAT or without.

Both computers used in the test system have a similar hardware and software:

- \rightarrow Intel Core i7 CPU
- → Intel PRO/1000 nic
- \rightarrow Ubuntu 12.04 / Kernel 3.8.0

The measurement was run using iperf 2.05. The TCP window size was set to 256 kb and five sessions were run simultaneously.

vRouter measurements were done with the following hardware:

- → PCs: Core i7-6700, Intel X540 10GE interfaces
- → ESXi Server: Dell PowerEdge R330 with Xeon E3-1230v5, 3.4 GHz, Intel X710 10GE interfaces as uplink to the ESXi vSwitches
- \rightarrow vRouter: VMXNET3 with virtual interfaces

IPerf was used without window parameter for the vRouter tests

IPSec routing performance

Other than with pure routing performance, IPSec-VPN routing actually changes the frames which are being passed from one interface to the next. When data is encrypted for the VPN tunnel, the original frame is encapsulated and it is supplemented with additional information. This has two important effects when considering the performance of IPSec routing:

→ Encrypted frames are larger than unencrypted frames. Consequently, any results have to indicate which frame size was observed at which interface, and/



or whether frames were encrypted or unencrypted. The values presented here always relate to an unencrypted frame size. An IP packet of 46 bytes is transported unencrypted, e.g. in a frame of 64 bytes. In the event of AES encryption, the frame grows for example to 122 bytes (46 byte IP packet + 14 byte Ethernet + 4 byte FCS + 20 byte IP + 8 byte ESP + 16 byte initialization vector (IV) + 0 byte padding (0-15 byte) + 1 byte padding length + 1 byte next header + 12 byte authentication).

→ The processes of encryption and decryption in the router take up computing time. These processes take place in two steps which, in the case of encryption, must be sequential. With decryption, on the other hand, these steps can be executed in parallel. Router models with VPN hardware acceleration provide significantly better performance with decryption than with encryption. This explains why the results display a significant difference in performance between the decryption and encryption directions. All of the IPSec-routing values given here are for a single VPN tunnel. With up to 1,000 tunnels established under laboratory conditions, the frame rate remained almost constant over all of the active tunnels. However, under actual operating conditions an increasing number of tunnels will cause the frame rate to drop due to the processes running for each tunnel (for example renewal of the key being used).

IPSec routing with different IMIX (decryption and encryption)

As an alternative to measurements with fixed frame sizes, series of measurements were performed with different IMIX patterns. The IMIX patterns simulate "real" data traffic composed of different frame sizes. There is no binding guideline for the composition of the frame sizes used, so in addition to the default setting of the IXIA test system (IMIX 0), two other common patterns were used for the measurement (IMIX 1 and IMIX 2). The individual patterns use the following frame compositions: \rightarrow IMIX 0: 45% 64 bytes, 20% 128 bytes, 5% 256 bytes, 3% 512 bytes, 2%

- 1024 bytes, 1% 1280 bytes, 24% 1364 bytes.
- \rightarrow IMIX 1: 7× 64 bytes, 4× 570 bytes, 1× 1418 bytes.
- → IMIX 2: 58% 90 bytes, 2% 92 bytes, 24% 594 bytes, 16% 1418 bytes.

Assuming an overhead of 100 bytes, 1418 bytes is the maximum frame size that can be transmitted encrypted on the Ethernet (with a maximum frame size for IEEE 802.3 of 1518 bytes).

Once again, it can be seen in these measurements that the decryption of the data is usually faster than the encryption.





Bridging performance (L2TPv3 Ethernet tunnel)

The bridging performance (<u>table 10</u>) examines the maximum data throughput that can be achieved via a transparent Ethernet tunnel established over an L2TPv3 pseudowire. Here, IPerf is used to perform a measurement with five parallel TCP streams through the Ethernet tunnel.

The hardware described under "Routing performance (TCP)" was used for the vRouter measurements.

A WAN-WAN connection is used for the vRouter measurements. A LAN-LAN connection is used for the remaining measurements.



Measurement values for devices with 1 Gbps interfaces

Note: For large frames or TCP, the measured performance may not be determined by the performance of the device, but is limited by the Ethernet interfaces (1G or 10G).

Table 1 - WAN-LAN routing – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.50	61.0 119,000	122 119,000	238 116,000	473 115,000	933 113,000	982 95,800	984 81,000	Mbps fps
1800EF, 1800EF-5G, 1800EFW	10.60	124 242,000	238 232,000	444 217,000	831 202,000	977 119,000	982 95,800	984 81,000	Mbps fps
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.50	84.7 165,000	169 165,000	339 165,000	674 164,000	977 119,000	982 95,800	984 81,000	Mbps fps
ISG-1000 WLC-1000	10.50	92.2 180,000	184 180,000	368 179,000	736 179,000	981 119,000	985 96,100	987 81,200	Mbps fps

Table 2 - LAN-LAN routing – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.50	73.7 144,000	145 141,000	290 141,000	580 141,000	977 119,000	982 95,800	984 81,000	Mbps fps
1800EF, 1800EF-5G, 1800EFW	10.60	157 306,000	298 290,000	540 263,000	955 233,000	977 119,000	982 95,800	984 81,000	Mbps fps
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.50	99.1 193,000	198 193,000	398 194,000	793 193,000	977 119,000	982 95,800	984 81,000	Mbps fps
ISG-1000 WLC-1000	10.50	125 245,000	249 243,000	504 246,000	962 234,000	981 119,000	985 96,100	987 81,200	Mbps fps



Table 3 - WAN-LAN routing – throughput [Mbps]

LCOS	Throughput with 5 TCP sessions
10.50	833
10.60	938
10.50	909
	10.50

Table 4 - TCP throughput [Mbps], simultaneous measurement in both directions
--

LANCOM	LCOS	Method	Throughput with 1 TCP session
1800EF, 1800EF-5G, 1800EFW	10.60	WAN-LAN routing	1,370
1800EF, 1800EF-5G, 1800EFW	10.60	LAN-LAN routing	1,570

Table 5 - IPSec routing AES-CBC UDP - decryption - throughput [Mbps] at frame size [bytes] and frame rate [fps]

				-	-	-	-		
LANCOM	LCOS	64	128	256	512	1024	1280	1418	
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.50	25.5 49,700	52.1 50,800	102 50,000	203 49,600	400 48,800	494 48,200	535 47,200	Mbps fps
1800EF, 1800EF-5G, 1800EFW	10.60	50.4 98,300	99.5 97,100	192 93,900	387 94,500	758 92,500	914 89,200	921 81,100	Mbps fps
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.50	40.3 78,600	80.5 78,500	161 78,600	316 77,100	638 77,800	794 77,500	875 77,100	Mbps fps
ISG-1000 WLC-1000	10.50	37.2 72,600	74.8 73,000	149 72,700	299 73,000	595 72,600	743 72,500	821 72,300	Mbps fps



Table 6 - IPSec routing AES-CBC - UDP - encryption - throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.50	29.4 57,500	58.6 57,100	116 56,800	230 56,100	449 54,700	550 53,700	606 53,400	Mbps fps
1800EF, 1800EF-5G, 1800EFW	10.60	59.0 115,000	117 114,000	229 111,000	458 111,000	858 104,000	942 92,000	948 83,500	Mbps fps
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.50	43.9 85,500	89.2 87,100	178 86,900	353 86,200	700 85,400	873 85,200	945 83,300	Mbps fps
ISG-1000 WLC-1000	10.50	42.0 82,000	85.7 83,700	171 83,700	343 83,800	684 83,500	853 83,200	944 83,200	Mbps fps

Table 7 / 8 - IPSec routing decryption / encryption – throughput [Mbps]

	1000	Decryptic	n		Encryptic	n	
LANCOM	LCOS	IMIX 0	IMIX 1	IMIX 2	IMIX 0	IMIX 1	IMIX 2
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.50	174	135	165	202	158	194
1800EF, 1800EF-5G, 1800EFW	10.60	343	275	323	413	326	389
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.50	261	204	249	306	237	292
ISG-1000 WLC-1000	10.50	252	197	241	296	230	282

Table 9 - TCP throughput [Mbps]

LANCOM	LCOS	Method	Throughput with 5 TCP sessions
ISG-1000	10.50	LAN - WAN (IPoE) routing	940
ISG-1000	10.50	LAN - LAN routing	940

Table 10 - L2TPv3 bridging – throughput [Mbps]
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LANCOM	LCOS	Throughput with 5 TCP sessions
ISG-1000 WLC-1000	10.50	991



Measurement values for devices with 10 Gbps interfaces

Note: For large frames or TCP, the measured performance may not be determined by the performance of the device, but is limited by the Ethernet interfaces (1G or 10G).

Table 11 - IPSec AES256-GCM - UDP decryption - throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
ISG-4000	10.50	44.9 87,600	93.8 91,600	189 92,300	380 92,700	754 92,000	940 91,700	1,040 92,000	Mbps fps
ISG-5000	10.60	221 431,000	448 437,000	894 436,000	1,740 426,000	3,260 398,000	3,830 373,000	4,030 355,000	Mbps fps
ISG-8000 Multicore	10.50	551 1,070,000	1,100 1,070,000	2,170 1,060,000	4,210 1,020,000	8,120 992,000	9,450 923,000	9,470 835,000	Mbps fps
vRouter	10.50	399 778,000	817 797,000	1,520 745,000	2,860 698,000	4,990 610,000	5,870 573,000	6,250 551,000	Mbps fps

Table 12 - IPSec AES256-GCM - UDP encryption - throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
ISG-4000	10.50	58.7 107,000	114 107,000	224 107,000	443 107,000	882 107,000	1,090 106,000	1,210 106,000	Mbps fps
ISG-5000	10.60	221 406,000	486 459,000	1,030 492,000	1,780 429,000	3,080 373,000	3,640 352,000	3,850 336,000	Mbps fps
ISG-8000 Multicore	10.50	560 1,020,000	1,100 1,040,000	2,140 1,020,000	4,170 1,000,000	7,680 928,000	9,540 923,000	9,500 835,000	Mbps fps
vRouter	10.50	357 656,000	602 570,000	1,280 618,000	2,490 603,000	4,170 508,000	4,670 455,000	4,660 409,000	Mbps fps

Table 13 - IPSec AES256-CBC SHA256 - UDP decryption - throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
ISG-4000	10.50	44.9 87,600	90.8 88,600	185 90,100	369 90,000	736 89,800	911 88,900	1,010 89,300	Mbps fps
ISG-5000	10.60	83.3 162,000	155 151,000	270 131,000	422 103,000	593 72,300	645 63,000	672 59,200	Mbps fps
ISG-8000 Multicore	10.50	186 362,000	353 334,000	632 308,000	1,060 259,000	1,570 192,000	1,740 170,000	1,830 161,000	Mbps fps
vRouter	10.50	220 429,000	415 405,000	723 353,000	1,000 244,000	1,760 215,000	1,840 179,000	2,040 180,000	Mbps fps

Table 14 - IPSec AES256-CBC SHA256 - UDP encryption - throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
ISG-4000	10.50	58.7 107,000	114 107,000	225 107,000	446 108,000	886 107,000	1,100 107,000	1,220 107,000	Mbps fps
ISG-5000	10.60	71.2 130,000	134 126,000	229 109,000	327 79,200	493 59,500	545 52,700	559 49,100	Mbps fps
ISG-8000 Multicore	10.50	170 311,000	314 294,000	547 263,000	891 215,000	1,280 155,000	1,410 137,000	1,480 130,000	Mbps fps
vRouter	10.50	188 345,000	301 285,000	617 296,000	927 224,000	1,440 173,000	1,580 153,000	1,640 144,000	Mbps fps



Table 15 / 16 - IPSec IMIX AES256-GCM - UDI	P decryption / encryption – throughp	out [Mbps]
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LANCOM	1005	Decryptic	on		Encryption		
LANCOM	LCOS	IMIX 0	IMIX 1	IMIX 2	IMIX 0	IMIX 1	IMIX 2
ISG-4000	10.50	335	252	303	397	295	356
ISG-5000	10.60	1,790	1,380	1,650	1,680	1,280	1,510
ISG-8000 Multicore	10.50	3,820	2,890	3,210	3,790	2,840	2,730
vRouter	10.50	2,600	2,050	2,520	1,680	1,440	1,840

Table 17 / 18 - IPSec IMIX AES256-CBC SHA256 - UDP decryption / encryption - throughput [Mbps]

	1000	Decryptic	on		Encryption		
LANCOM	LCOS	IMIX 0	IMIX 1	IMIX 2	IMIX 0	IMIX 1	IMIX 2
ISG-4000	10.50	333	252	300	396	294	358
ISG-5000	10.60	400	330	382	319	264	310
ISG-8000 Multicore	10.50	985	795	931	830	681	787
vRouter	10.50	1,140	924	1,060	926	771	732

Table 19 - TCP throughput [Mbps]

LANCOM	LCOS Method	Throughput 5 TCP sessions
ISG-4000	10.50 LAN - LAN	3,542
ISG-4000	10.50 LAN - WAN (IPoE)	2,069
ISG-4000	10.50 WAN - LAN (IPoE)	2,201
ISG-5000	10.60 LAN - LAN	7,695
ISG-5000	10.60 LAN - WAN (IPoE)	4,968
ISG-5000	10.60 WAN - LAN (IPoE)	5,101
ISG-8000 Multicore	10.50 LAN - LAN	9,400
ISG-8000 Multicore	10.50 LAN - WAN (IPoE)	9,400
ISG-8000 Multicore	10.50 WAN - LAN (IPoE)	9,400
vRouter	10.50 LAN - LAN	9,400
vRouter	10.50 LAN - WAN (IPoE)	9,400
vRouter	10.50 WAN - LAN (IPoE)	9,400



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