

Routing performance LCOS 10.70

Applications for communications and entertainment are increasingly based on IP networks. In order to ensure that the necessary bandwidth performance can be provided reliably, it is important for the infrastructure's networking components to be tested thoroughly and intensively. In this techpaper, LANCOM Systems presents the methods of measuring routing and VPN performance for central site and VPN gateways as well as the respective results.

We have examined a variety of aspects for consideration when measuring the router performance. This includes transmission speeds of connections between the LAN and the Internet (WAN), and the internal data transmission in the network (LAN-LAN). Many business processes rely on secure WAN connections, which is why we have focused on determining the performance of encrypted data connections over VPN.

Test system

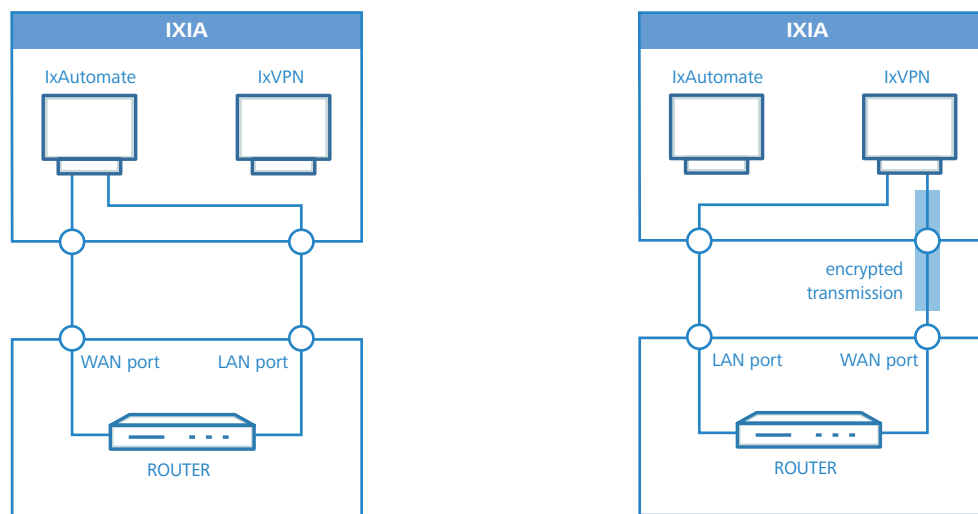


Figure 1:
IXIA test system for
routing connections
and encrypted VPN
connections between LAN
and WAN

All of the performance values were measured in the LANCOM test laboratory. Tests were conducted with an IXIA test system. IXIA uses so-called test suites, which enable the simulation of different applications. This allows, for example, the investigation of data throughput over automatically established VPN tunnels, or the measurement of pure LAN-WAN routing performance for unidirectional and bidirectional data connections. IXIA is a leading supplier of systems which test IP-based services and infrastructures. Test systems from IXIA are employed all over the world by network component manufacturers and other organizations to help assure the functionality and reliability of complex IP networks, devices, and applications.

The measurement of data transmission itself uses either a fixed frame size or a combination of frame sizes which reflects a typical flow of data. These combinations are known as "Internet Mix", or IMIX for short. The type of IMIX which is applied significantly affects the test results because packet size has a strong influence on a connection's performance. By selecting the appropriate ports on the router being tested, it is possible to test connections between the LAN and the WAN, and also pure LAN-LAN connections.

The setup for measuring transfer rates > 1 GBit/s represents a large central-site scenario. In this scenario, several central sites can also work as a network, which is why an intermediate router with BGP ensures that the packets for each tunnel pass through the respective central site (measured values from [Table 8](#)).

Two scenarios are used for the vRouter measurements. One is the transmission from a PC in the LAN to a PC in another LAN or WAN by using the vRouter in between. The second scenario adds two vRouters which are connected over WAN and use a VPN tunnel for encrypted data transfer.

Firmware on the devices under test is LCOS 10.70 (unless otherwise stated).

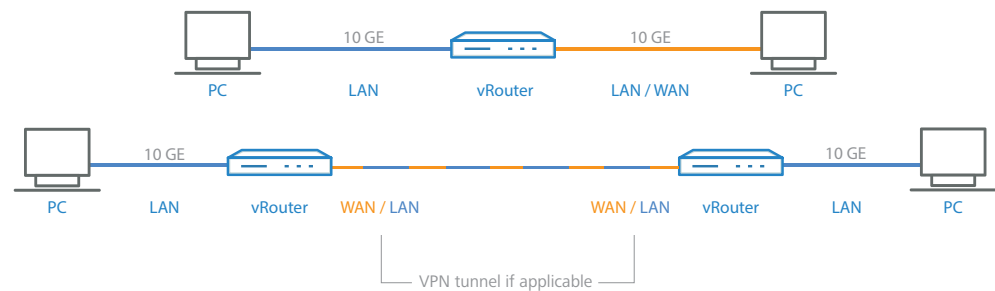


Figure 2:
Schematic view of the
vRouter test system
scenarios

Routing performance (UDP)

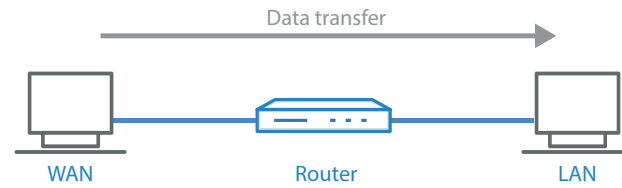


Figure 3:
Schematic view of the
test system

The measurement of routing performance involves the determination of the maximum data throughput which can be achieved before a router starts rejecting packets. This measurement uses UDP packets of various sizes in order to simulate the performance with different applications. Ethernet frame sizes range from 64 bytes for the smallest to 1518 bytes for the largest frames. Tests on different router models demonstrate the influence of the different hardware platforms (processor, interfaces).

Measurements initially determine the frame rate, which is a good performance indicator of the tested hardware. With normal routing, the frame rate is fairly constant even with different frame sizes. This is because only the header is inspected during routing, a process which is largely independent of the size of the frames being routed. For this reason, only the typical frame rates are given in the tables.

The throughput for a certain frame size (or even a mix of sizes, see [IMIX](#) on page 05) can be approximately calculated by multiplication with the frame rate. When the frame rate is constant, data throughput depends directly on the frame size because the larger the frames, the larger is the data volume that can be transmitted. The maximum number of frames transmitted per second is limited by the performance of the interfaces and the transmission medium.

Measurement of the routing performance relates to the size of the Ethernet frames. To compare packet sizes for particular applications, it is necessary to subtract the header. For a frame of 512 bytes, the result is a UDP datagram size of 474 bytes (512 bytes - 14 bytes Ethernet header - 4 bytes FCS trailer - 20 bytes IP header) and, after subtracting the UDP header (8 bytes), the UDP payload is 466 bytes.

To investigate routing performance, in this paper two different applications are considered.

→ For WAN-LAN routing, data received from the WAN is forwarded to a peer in the LAN.

→ For LAN-LAN routing, data remains within the local-area network and is passed from one LAN port to another.

The measurements show that the throughput increases almost linearly with the frame size until the limit of the Gigabit interface is reached.

Routing performance (TCP)

Using UDP, the maximum performance is easily measured. However, a large portion of data transfers utilizes TCP. Therefore, it is important to take a closer look at those scenarios as well.

The TCP measurement is accomplished by using iperf, a program used to measure the data throughput between two computers. They are connected to a LANCOM router. One computer acts as a server in the WAN, sending data packets to the other, which acts as a client in the LAN, representing a typical download scenario.

The TCP measurement done with iperf did not use NAT since no direct measurement with NAT is possible using iperf. The correlation between the test results and the results seen at a WAN connection including NAT in a live environment is very high. The causes are the differences of the traffic structure of iperf compared to an HTTP download and the LCOS firewall which offers a similar performance with NAT or without.

Both computers used in the test system have a similar hardware and software:

- Intel Core i7 CPU
- Intel PRO/1000 nic
- Ubuntu 12.04 / Kernel 3.8.0

The measurement was run using iperf 2.05. The TCP window size was set to 256 kb and five sessions were run simultaneously.

vRouter measurements were done with the following hardware:

- PCs: Core i7-6700, Intel X540 10GE interfaces
- ESXi Server: Dell PowerEdge R330 with Xeon E3-1230v5, 3.4 GHz, Intel X710 10GE interfaces as uplink to the ESXi vSwitches
- vRouter: VMXNET3 with virtual interfaces

IPerf was used without window parameter for the vRouter tests

IPSec routing performance

Other than with pure routing performance, IPSec-VPN routing actually changes the frames which are being passed from one interface to the next. When data is encrypted for the VPN tunnel, the original frame is encapsulated and it is supplemented with additional information. This has two important effects when considering the performance of IPSec routing:

- Encrypted frames are larger than unencrypted frames. Consequently, any results have to indicate which frame size was observed at which interface, and/

or whether frames were encrypted or unencrypted. The values presented here always relate to an unencrypted frame size. An IP packet of 46 bytes is transported unencrypted, e.g. in a frame of 64 bytes. In the event of AES encryption, the frame grows for example to 122 bytes (46 byte IP packet + 14 byte Ethernet + 4 byte FCS + 20 byte IP + 8 byte ESP + 16 byte initialization vector (IV) + 0 byte padding (0-15 byte) + 1 byte padding length + 1 byte next header + 12 byte authentication).

→ The processes of encryption and decryption in the router take up computing time. These processes take place in two steps which, in the case of encryption, must be sequential. With decryption, on the other hand, these steps can be executed in parallel. Router models with VPN hardware acceleration provide significantly better performance with decryption than with encryption. This explains why the results display a significant difference in performance between the decryption and encryption directions. All of the IPSec-routing values given here are for a single VPN tunnel. With up to 1,000 tunnels established under laboratory conditions, the frame rate remained almost constant over all of the active tunnels. However, under actual operating conditions an increasing number of tunnels will cause the frame rate to drop due to the processes running for each tunnel (for example renewal of the key being used).

IPSec routing with different IMIX (decryption and encryption)

As an alternative to measurements with fixed frame sizes, series of measurements were performed with different IMIX patterns. The IMIX patterns simulate „real“ data traffic composed of different frame sizes. There is no binding guideline for the composition of the frame sizes used, so in addition to the default setting of the IXIA test system (IMIX 0), two other common patterns were used for the measurement (IMIX 1 and IMIX 2). The individual patterns use the following frame compositions:

- IMIX 0: 45% 64 bytes, 20% 128 bytes, 5% 256 bytes, 3% 512 bytes, 2% 1024 bytes, 1% 1280 bytes, 24% 1364 bytes.
- IMIX 1: 7× 64 bytes, 4× 570 bytes, 1× 1418 bytes.
- IMIX 2: 58% 90 bytes, 2% 92 bytes, 24% 594 bytes, 16% 1418 bytes.

Assuming an overhead of 100 bytes, 1418 bytes is the maximum frame size that can be transmitted encrypted on the Ethernet (with a maximum frame size for IEEE 802.3 of 1518 bytes).

Once again, it can be seen in these measurements that the decryption of the data is usually faster than the encryption.

Measurement values for devices with 1 Gbps interfaces

Note: For large frames or TCP, the measured performance may not be determined by the performance of the device, but is limited by the Ethernet interfaces (1G or 10G).

Table 1 - WAN-LAN routing – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1518	
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.70	58.7 114,000	117 114,000	229 112,000	444 108,000	889 108,000	982 95,800	984 81,000	MBit/s Frames/s
1800EF, 1800EF-5G, 1800EFW	10.70	121 235,000	234 228,000	440 214,000	831 202,000	977 119,000	982 95,800	984 81,000	MBit/s Frames/s
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.70	94.7 184,000	189 184,000	380 185,000	755 184,000	977 119,000	982 95,800	984 81,000	MBit/s Frames/s
WLC-1000	10.70	74.1 144,000	149 145,000	298 145,000	594 145,000	981 119,000	985 96,100	987 81,200	MBit/s Frames/s

Table 2 - LAN-LAN routing – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1518	
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.70	67.7 132,000	129 125,000	269 131,000	528 128,000	964 117,000	982 95,800	984 81,000	MBit/s Frames/s
1800EF, 1800EF-5G, 1800EFW	10.70	143 279,000	275 268,000	510 249,000	955 233,000	977 119,000	982 95,800	984 81,000	MBit/s Frames/s
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.70	107 209,000	216 211,000	432 211,000	859 209,000	977 119,000	982 95,800	984 81,000	MBit/s Frames/s
WLC-1000	10.70	104 202,000	207 202,000	414 202,000	831 202,000	981 119,000	985 96,100	987 81,200	MBit/s Frames/s

Table 3 - WAN-LAN routing – throughput [Mbps]

LANCOM	LCOS	Throughput with 5 TCP sessions
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.70	765
1800EF, 1800EF-5G, 1800EFW	10.70	938
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.70	909

Table 4 - TCP throughput [Mbps], simultaneous measurement in both directions

LANCOM	LCOS	Method	Throughput with 1 TCP session
1800EF, 1800EF-5G, 1800EFW	10.70	WAN-LAN routing	1,370
1800EF, 1800EF-5G, 1800EFW	10.70	LAN-LAN routing	1,570

Table 5 - IPsec routing AES-CBC UDP - decryption – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.70	38.2 74,500	76.3 74,500	151 73,500	299 73,000	567 69,100	689 67,300	764 67,300	MBit/s Frames/s
1800EF, 1800EF-5G, 1800EFW	10.70	102 198,000	199 194,000	385 188,000	773 188,000	918 112,000	928 90,500	928 81,800	MBit/s Frames/s
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.70	62.7 122,000	124 121,000	250 121,000	498 121,000	920 112,000	925 90,300	905 79,800	MBit/s Frames/s
WLC-1000	10.70	61.3 119,000	121 117,000	241 117,000	485 118,000	921 112,000	928 90,500	928 81,800	MBit/s Frames/s

Table 6 - IPSec routing AES-CBC - UDP - encryption – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.70	36.8 71,800	72.8 71,100	144 70,500	286 69,800	545 66,500	669 65,300	735 64,800	MBit/s Frames/s
1800EF, 1800EF-5G, 1800EFW	10.70	91.7 179,000	180 175,000	351 171,000	681 166,000	914 111,000	937 91,400	947 83,400	MBit/s Frames/s
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.70	62.9 122,000	127 123,000	253 123,000	503 122,000	926 113,000	940 91,700	945 83,300	MBit/s Frames/s
WLC-1000	10.70	57.1 111,000	114 111,000	228 111,000	461 112,000	916 111,000	942 92,000	948 83,500	MBit/s Frames/s

Table 7 - IPSec routing decryption / encryption – throughput [Mbps]

LANCOM	LCOS	Decryption			Encryption		
		IMIX 0	IMIX 1	IMIX 2	IMIX 0	IMIX 1	IMIX 2
730VA, 883 VoIP, 884 VoIP 1640E 1780EW-4G+ 1781EW+, 1781VA, 1781VAW 1783VA, 1783VAW, 1784VA 1790-4G, 1790EF, 1790VA, 1790VA-4G, 1790VAW 1793VA, 1793VA-4G, 1793VAW IAP-1781VAW+ IAP-4G+	10.70	258	201	244	251	195	236
1800EF, 1800EF-5G, 1800EFW	10.70	688	544	652	572	455	560
1900EF, 1900EF-5G 1906VA, 1906VA-4G 1926VAG, 1926VAG-4G, 1926VAG-5G	10.70	414	320	393	434	336	413
WLC-1000	10.70	395	310	380	398	308	376

Measurement values for devices with 10 Gbps interfaces

Note: For large frames or TCP, the measured performance may not be determined by the performance of the device, but is limited by the Ethernet interfaces (1G or 10G).

Table 8 - IPSec AES256-GCM - UDP decryption – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
ISG-5000	10.70	290	574	1,120	2,040	3,520	4,050	4,220	MBit/s
		566,000	560,000	547,000	500,000	430,000	396,000	371,000	Frames/s
ISG-8000	10.70	655	1,310	2,620	5,270	9,320	9,450	9,480	MBit/s
		1,270,000	1,280,000	1,280,000	1,280,000	1,130,000	923,000	835,000	Frames/s
vRouter	10.70	412	823	1,600	2,790	5,490	5,960	6,180	MBit/s
		805,000	803,000	785,000	682,000	671,000	582,000	545,000	Frames/s

Table 9 - IPSec AES256-GCM - UDP encryption – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
ISG-5000	10.70	280	543	1,060	1,930	3,270	3,870	4,110	MBit/s
		511,000	511,000	510,000	466,000	394,000	375,000	360,000	Frames/s
ISG-8000	10.70	611	1,190	2,460	4,970	9,450	9,540	9,500	MBit/s
		1,120,000	1,110,000	1,180,000	1,190,000	1,130,000	923,000	835,000	Frames/s
vRouter	10.70	350	764	1,430	2,720	4,370	5,250	5,580	MBit/s
		642,000	723,000	687,000	660,000	532,000	511,000	489,000	Frames/s

Table 10 - IPSec AES256-CBC SHA256 - UDP decryption – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
ISG-5000	10.70	172	441	714	1,000	1,320	1,420	1,470	MBit/s
		335,000	430,000	348,000	245,000	162,000	138,000	129,000	Frames/s
ISG-8000	10.70	442	868	1,690	3,320	4,330	4,630	4,780	MBit/s
		863,000	847,000	828,000	811,000	529,000	452,000	421,000	Frames/s
vRouter	10.70	308	575	968	1,540	2,090	2,260	2,350	MBit/s
		600,000	561,000	472,000	376,000	255,000	220,000	207,000	Frames/s

Table 11 - IPSec AES256-CBC SHA256 - UDP encryption – throughput [Mbps] at frame size [bytes] and frame rate [fps]

LANCOM	LCOS	64	128	256	512	1024	1280	1418	
ISG-5000	10.70	81.2	154	258	386	528	579	595	MBit/s
		149,000	144,000	123,000	93,500	64,200	55,900	52,300	Frames/s
ISG-8000	10.70	206	370	641	1,000	1,390	1,520	1,590	MBit/s
		379,000	350,000	305,000	242,000	169,000	148,000	140,000	Frames/s
vRouter	10.70	210	384	668	1,060	1,430	1,610	1,710	MBit/s
		385,000	364,000	321,000	255,000	173,000	156,000	150,000	Frames/s

Table 12 - IPSec IMIX AES256-GCM - UDP decryption / encryption – throughput [Mbps]

LANCOM	LCOS	Decryption			Encryption		
		IMIX 0	IMIX 1	IMIX 2	IMIX 0	IMIX 1	IMIX 2
ISG-5000	10.70	1,760	1,470	1,640	1,810	1,430	1,590
ISG-8000	10.70	4,680	3,480	4,250	4,340	3,200	4,050
vRouter	10.70	2,490	1,960	2,520	2,080	1,440	1,490

Table 13 - IPSec IMIX AES256-CBC SHA256 - UDP decryption / encryption – throughput [Mbps]

LANCOM	LCOS	Decryption			Encryption		
		IMIX 0	IMIX 1	IMIX 2	IMIX 0	IMIX 1	IMIX 2
ISG-5000	10.70	997	849	963	371	312	356
ISG-8000	10.70	3,030	2,280	2,720	941	777	894
vRouter	10.70	1,430	1,240	1,310	834	819	901

Table 14 - TCP throughput [Mbps]

LANCOM	LCOS	Method	Throughput 5 TCP sessions
ISG-5000	10.70	LAN - LAN	7,000
ISG-5000	10.70	LAN - WAN (IPoE)	4,910
ISG-5000	10.70	WAN - LAN (IPoE)	5,320
ISG-8000	10.70	LAN - LAN	9,400
ISG-8000	10.70	LAN - WAN (IPoE)	9,400
ISG-8000	10.70	WAN - LAN (IPoE)	9,400
vRouter	10.70	LAN - LAN	9,400
vRouter	10.70	LAN - WAN (IPoE)	9,400
vRouter	10.70	WAN - LAN (IPoE)	9,400